

Metamodel



Run-to-Completion Step: Overview



- Choose an **event** from the event pool (queue)
- Choose a **maximal**, **conflict-free**, **prioritized**, set of transitions enabled by the event
- Execute set of transitions
 - exit source states (inside-out)
 - execute transition effects
 - enter target states (outside-in)

thereby generating new events and activities

- Active state configuration
 - the states the state machine currently is in
 - forms a tree
 - if a composite state is active, all its regions are active
- Least-common-ancestor (LCA) of states s_1 and s_2
 - the least region or orthogonal state (upwards) containing s_1 and s_2



bold: active state configuration



bold: LCA of states A1 and A2

- Compound transitions
 - transitions for an event are "chained" into compound transitions
 - eliminating pseudostates like junction, fork, join, entry, exit
 - this is not possible for choice pseudostates where the guard of outgoing transitions are evaluated dynamically (in contrast to junctions)
 - several source and target states





Run-to-Completion Step: Preliminaries (3)

- Main source / target state *m* of compound transition *t*
 - Let *s* be LCA of all source and target states of *t*
 - If s region: *m* = direct subvertex of *s* containing all source states of *t*
 - If *s* orthogonal state: *m* = *s*
 - Similarly for main target state
 - All states between main source and explicit source states are exited, all state between main target and explicit target states are entered.
- Conflict of compound transitions t₁ and t₂
 - intersection of states exited by t_1 and t_2 not empty
- **Priority** of compound transition t_1 over t_2
 - s_i "deepest" source state of transition t_i
 - s₁ (direct or transitive) substate of s₂

```
\mathsf{RTC}(env, conf) \equiv
  event \leftarrow fetch()
    step \leftarrow choose steps(conf, event)
    if step = \emptyset \land event \in deferred(conf)
    then defer(event)
    fi
    for transition \in step do
      conf \leftarrow handleTransition(env, conf, transition)
    od
    if isCall (event) \land event \notin deferred(conf)
    then acknowledge(event)
    fi
    conf ]
```



Run-to-Completion Step (2)

```
steps(env, conf, event) \equiv
   | transitions \leftarrow enabled(env, conf, event)
    {step | (guard, step) \in steps(conf, transitions) \land env \models guard }  \rfloor 
steps(conf, transitions) \equiv
   | steps \leftarrow \{(true, \emptyset)\}
    for transition \in transitions do
         for (guard, step) \in steps(conf, transitions \setminus \{transition\}) do
            if inConflict(conf, transition, step)
            then if higherPriority(conf, transition, step)
                  then guard \leftarrow guard \land \neg guard(transition) fi
            else step \leftarrow step \cup {transition}
                  guard \leftarrow guard \land guard(transition) fi
            steps \leftarrow steps \cup {(guard, step)} od od
    steps_
```

Run-to-Completion Step (3)

```
handleTransition(conf, transition) \equiv
  for state \in insideOut(exited(transition)) do
     uncomplete(state)
     for timer ∈ timers(state) do stopTimer(timer) od
      execute(exit(state))
     conf \leftarrow conf \setminus \{state\}
  od
  execute(effect(transition))
  for state \in outsideIn(entered(transition)) do
      execute(entry(state))
     for timer ∈ timers(state) do startTimer(timer) od
     conf \leftarrow conf \cup \{state\}
      complete(conf, state)
  od
  conf ]
```

Semantic variation points

- Some semantic variation points have been mentioned before.
 - delays in event pool
 - handling of deferred events
 - entering of composite states without default entry
- Which events are prioritized?
 - completion events only
 - all internal events (completion, time, change)
- Which (additional) timing assumptions?
 - delays in communication
 - time for run-to-completion step
 - zero-time assumption

State machine refinement

• State machines are behaviors and may thus be refined.





no refinement possible



Protocol state machines

- Protocol state machines specify which behavioral features of a classifier can be called in which state and under which condition and what effects are expected.
 - particularly useful for object life cycles and ports
 - no effects on transitions, only effect descriptions



Several operation specifications are combined conjunctively:

```
context C::op()
                                                            [P_1] op() / [Q_1]
   pre: inState(S_1) and P_1
                                                   S_1
   post: Q_1 and inState(S_3)
   context C::op()
                                                            [P_2]op()/[Q_2]
   pre: inState(S_2) and P_2
                                                   S_2
                                                                                S_{\Delta}
   post: Q_2 and inState(S_4)
results in
   context C::op()
   pre: (inState(S_1) and P_1) or (inState(S_2) and P_2)
   post: (inState@pre(S_1) and P_1@pre) implies (Q_1 and inState(S_3))
     and (inState@pre(S_2) and P_2@pre) implies (Q_2 and inState(S_4))
```

How things work together

- Static structure
 - sets the scene for state machine behavior
 - state machines refer to
 - properties
 - behavioral features (operations, receptions)
 - signals
- Interactions
 - may be used to exemplify the communication of state machines
 - refer to event occurrences used in state machines
- OCL
 - may be used to specify guards and pre-/post-conditions
 - refers to actions of state machines (OclMessage)
- Protocols and components
 - state machines may specify protocol roles

Wrap up

- State machines model behaviour
 - object and use case life cycles
 - control automata
 - protocols
- State machines consist of
 - Regions and ...
 - ... (Pseudo)States (with entry, exit, and do-activities) ...
 - connected by Transitions (with triggers, guards, and effects)
- State machines communicate via event pools.
- State machines are executed by run-to-completion steps.